

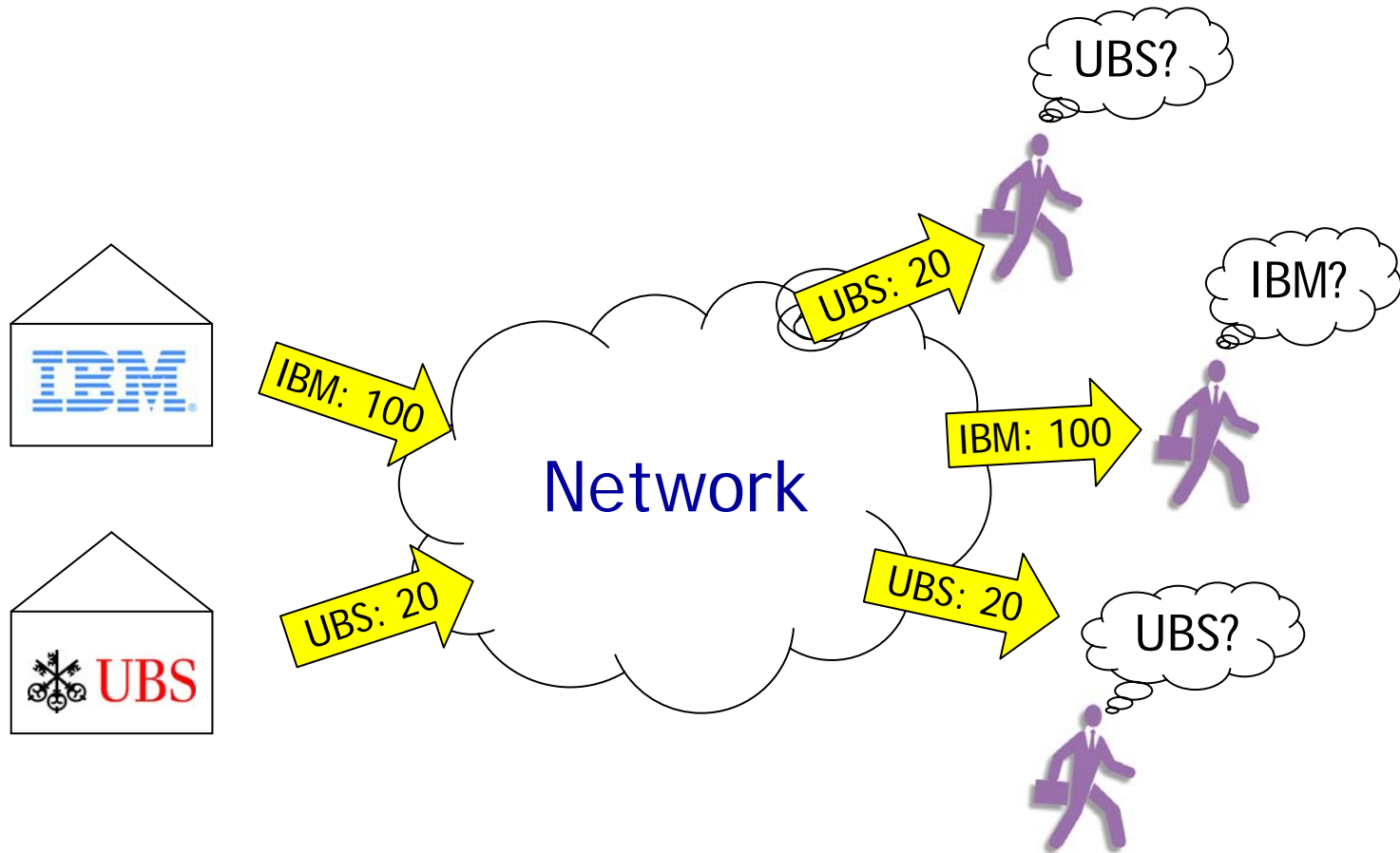


# **Publish/Subscribe Systems**

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# An Application Case



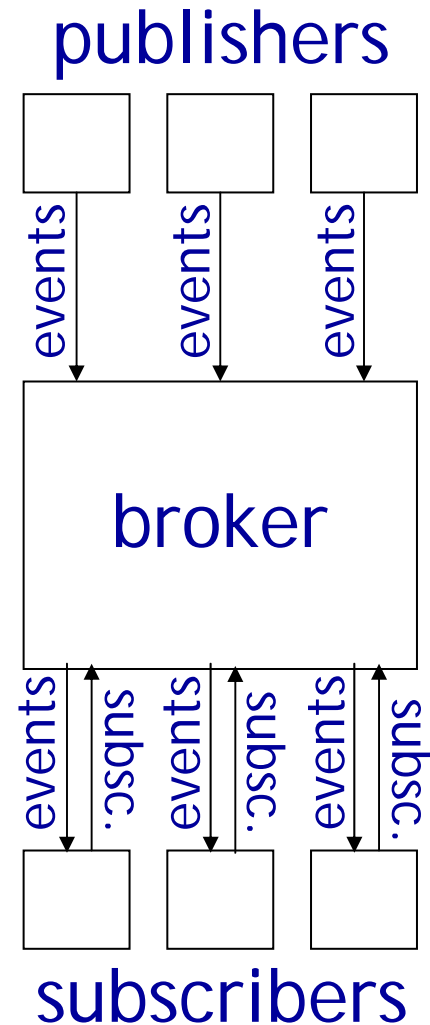
# Characteristics

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- Hard to realize with RPC
  - Sender needs to know identities of all receivers
  - Sender communicates with a large and dynamically changing set of receivers
  - Senders and receivers may not be up simultaneously
  - Information dissemination blocks sender
- Many other applications
  - Stock information systems, auction systems, air traffic control, news ticker, alerting services, ...

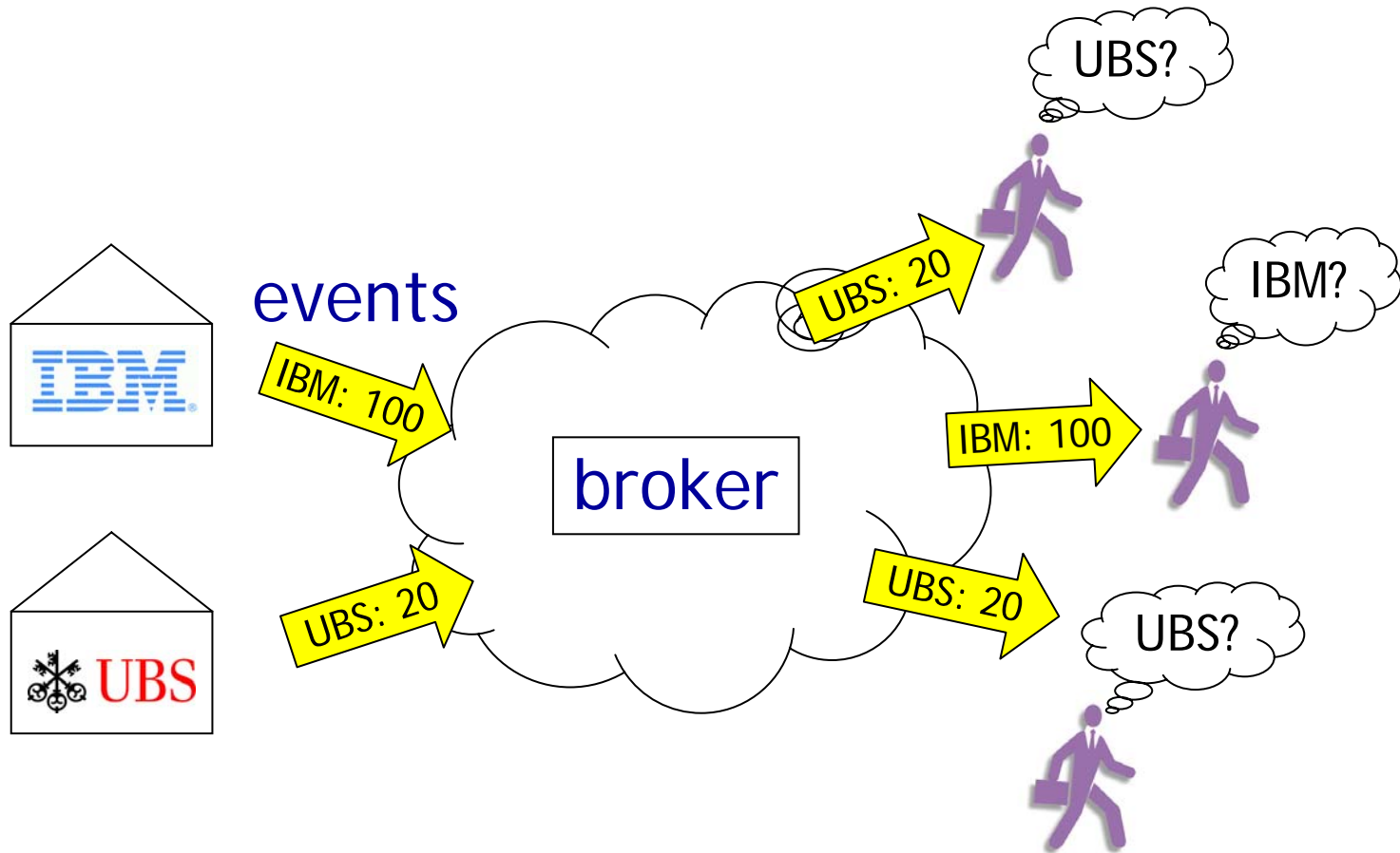
# Pub/Sub Architecture

- Publishers
  - Entities that produce information
- Subscribers
  - Entities that consume information
- Broker
  - Delivers information from producers to consumers
- Events
  - Basic units of information
- Subscription
  - Expression of interest in certain events



# Application Revisited

subscriptions



publishers

subscribers

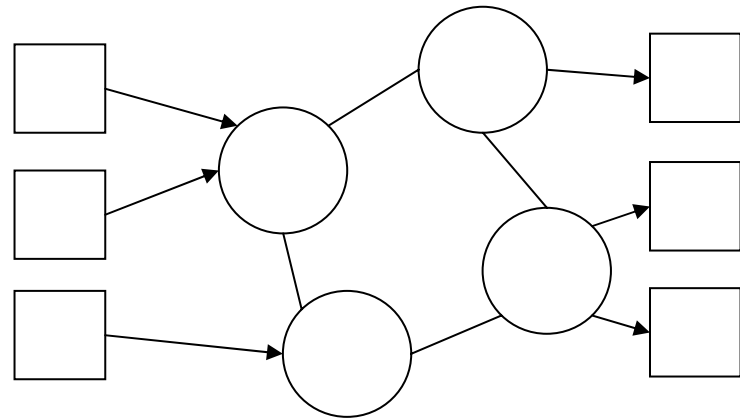
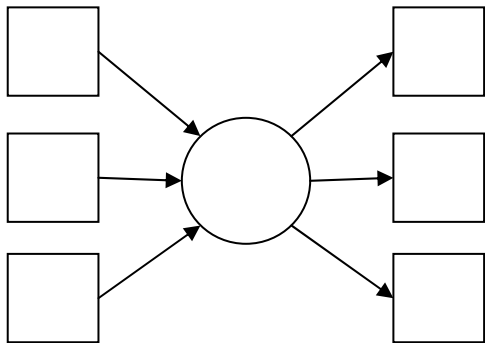
# Events & Notifications

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- Events
  - Asynchronous state transitions
  - Have a type: „stockvalue“
  - May have parameters: „IBM“ , „100“
    - Parameter names: „company“ , „value“
    - Parameter types: string, integer
    - Special parameters: time, location
- Representation
  - Plain text, XML, binary encoding, ...
  - Example: stockvalue(string company=„IBM“ , int value=100)
- Notifications
  - Messages carrying event representations

# Broker

- Main tasks
  - (De)registration of publishers / subscribers
  - Matching of events with subscriptions
  - Routing of events from publishers to matching subscribers
- Different implementations
  - Centralized vs. Distributed
  - Simplicity vs. Scalability



# Benefits

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- Decoupling of publishers and subscribers
  - Space: no need to be connected or even know each other (anonymous)
  - Time: no need to be up at the same time
  - Synchronization: publishing and receiving events are asynchronous operations
- Suitable for large, dynamic systems
  - Many participating entities
  - Entities frequently enter and leave the system



# Event Matching

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- Key operation to identify subscribers to which an event must be delivered
  - Input: event  $e$ , set of subscriptions  $S$
  - Output: all  $s \in S$  that match  $e$
- Depends on the semantics of subscriptions
  - Type-based: type of the event
  - Channel-based: cf. multicast group
  - Topic-based: cf. news group
  - Content-based: parameters of events

# Type-Based Subscriptions

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- Subscriptions specify the type of the requested events
  - E.g., „stockvalue“
- Event matching
  - Comparison of event type and suscription type

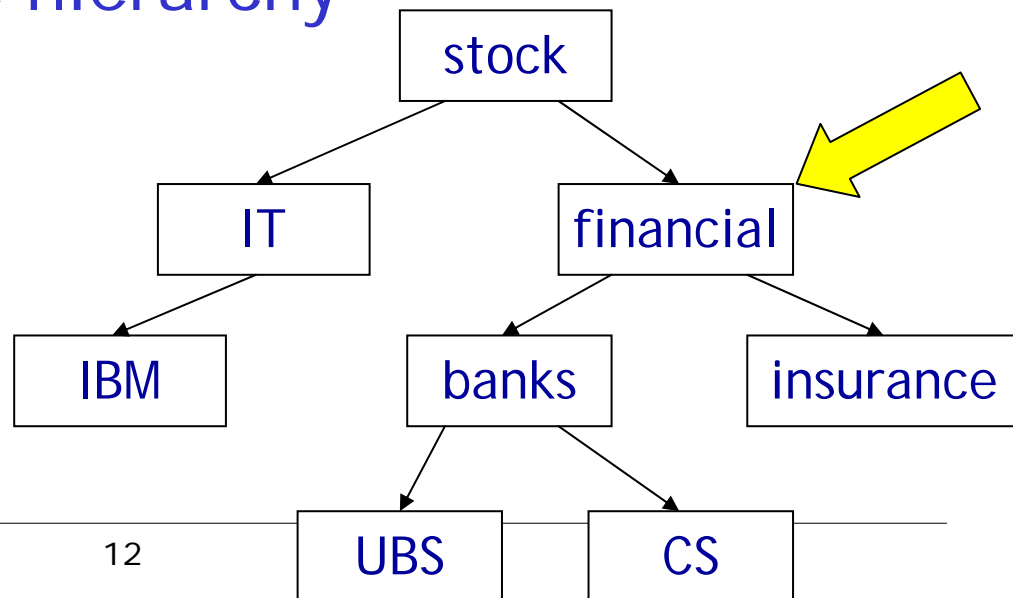
# Channel-based Subscriptions

- Multiple, named broadcasts channels can be created
  - E.g., news, stock, ...
- Publishers publish events to a specific channel
  - Events of different types may be published to a single channel
- Subscriptions specify the desired channel
- Event matching
  - E.g., a multicast group for each channel



# Topic-Based Subscriptions

- Each event has a „topic“ parameter
- Topics are organized in a hierarchy
- Subscriptions specify a topic of interest
- Event matching:
  - $s$  matches  $e$  if topic of  $e$  is descendent of topic of  $s$  in the hierarchy



# Content-Based Subscriptions

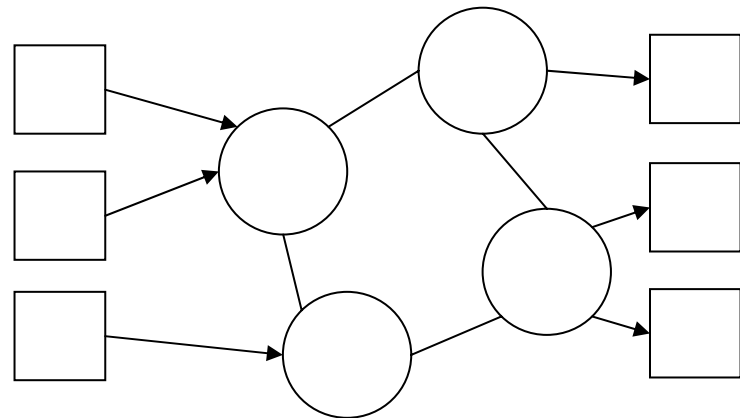
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- Subscriptions refer to type and parameter values of events
  - E.g.,  $type = „stockvalue“ \wedge company = „UBS“ \wedge value < 100$
  - Often a conjunction over atomic predicates:  
 $(p_1 \ op_1 \ v_1) \wedge (p_2 \ op_2 \ v_2) \wedge \dots$
- Event matching
  - $s$  matches  $e$  if the predicate of  $s$  holds on the parameter values of  $e$
  - Algorithms to process multiple subscriptions with much lower overhead
    - If „ $value < 100$ “ doesn't match, then „ $value < 20$ “ won't match as well

# Event Routing

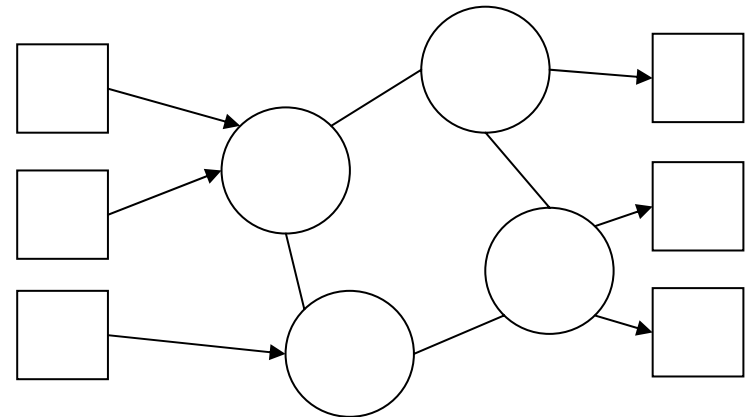
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- Nontrivial problem with distributed brokers for content-based subscriptions
  - Cannot use address-based routing!
- Many different approaches
  - Hierarchical vs. peer-to-peer
  - Flooding vs. overlay



# A Distributed Broker Model

- Subscribers and publishers connect to a local broker
- Brokers maintain connections to some, but not all other brokers
  - Form connected (acyclic?) network
- Routing problem:
  - Broker receives an event from either a publisher or another broker
  - To which brokers (and subscribers) should the event be sent?



# Event Flooding

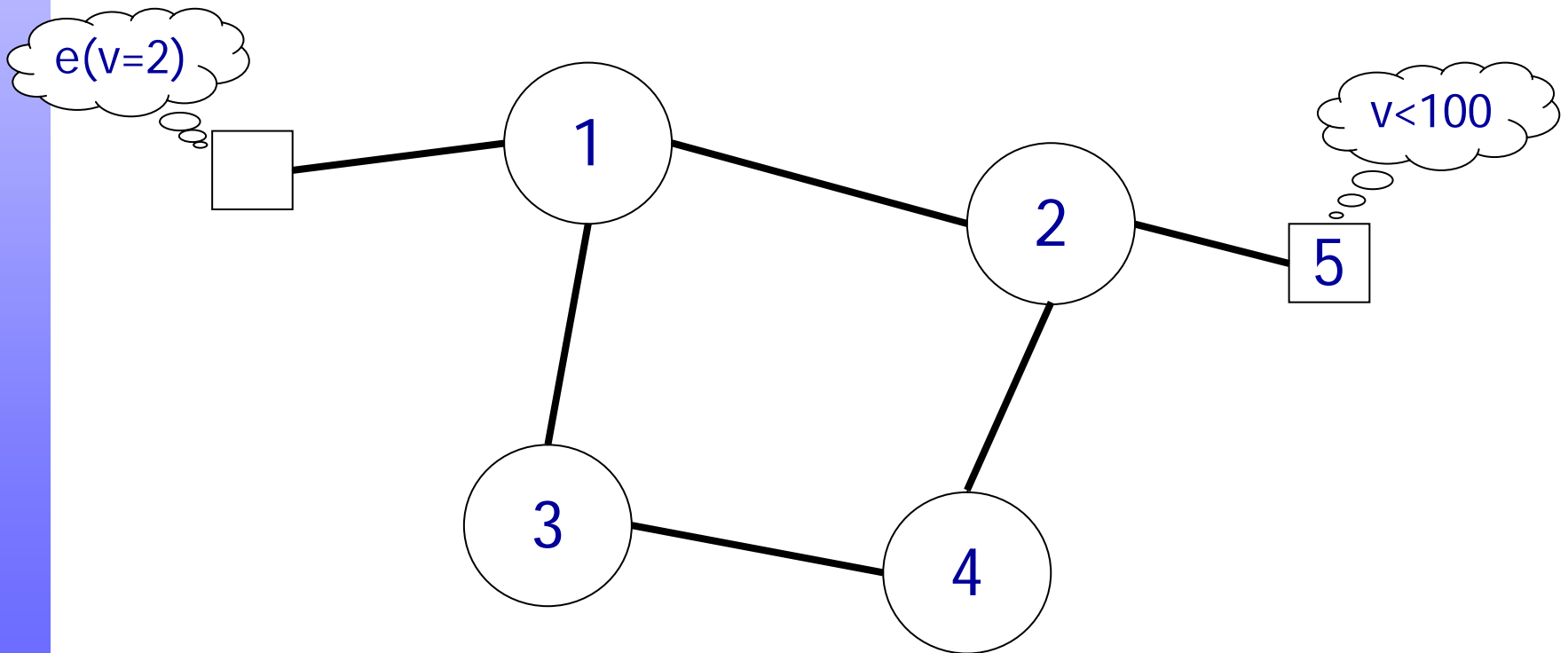
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- If a broker receives an event  $e$  from  $n$ 
  - Do nothing if  $e$  has been received before
  - Send  $e$  to all neighboring brokers except  $n$
  - Send  $e$  to all matching subscribers
- Ensures that all brokers will receive each event



# Example

- Event:  $e(v=2)$
- Content-based subscription:  $v < 100$



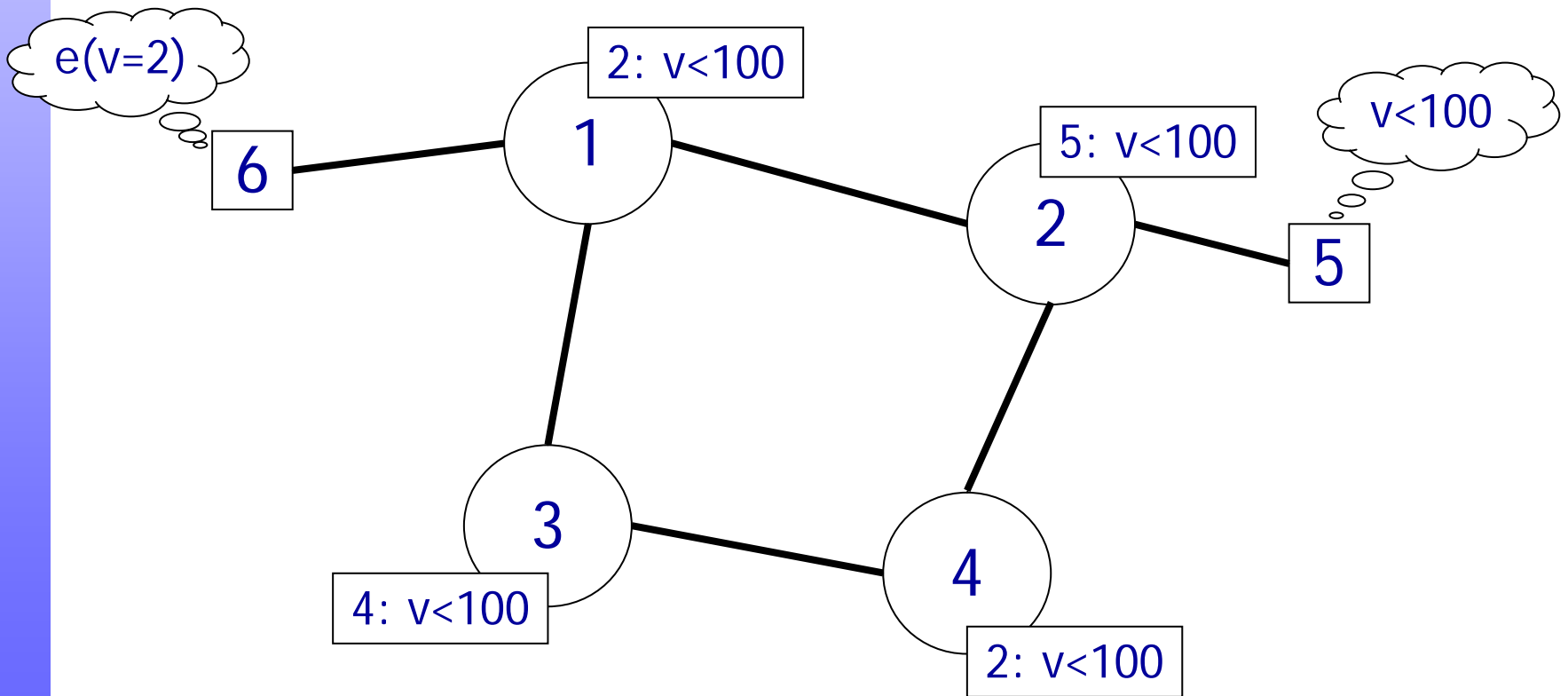
# Subscription Flooding

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- Idea: Use subscriptions to compute routing paths for events
  - This causes overhead whenever a subscriptions is added / removed / changed
  - But pays off if subscriptions change rarely compared to number of events published
- Approach: each broker maintains a routing table of entries  $(n_i, s_i)$ 
  - $n$  = neighboring broker,  $s$  = subscription
  - Send  $e$  to all  $n_i$  where  $s_i$  matches  $e$
- Setup by subscription flooding: When a broker receives  $s$  from  $n$ 
  - Do nothing if  $s$  has been received before
  - Create routing table entry  $(n, s)$
  - Send  $s$  to all neighboring brokers

# Example

- Event:  $e(v=2)$
- Content-based subscription:  $v < 100$



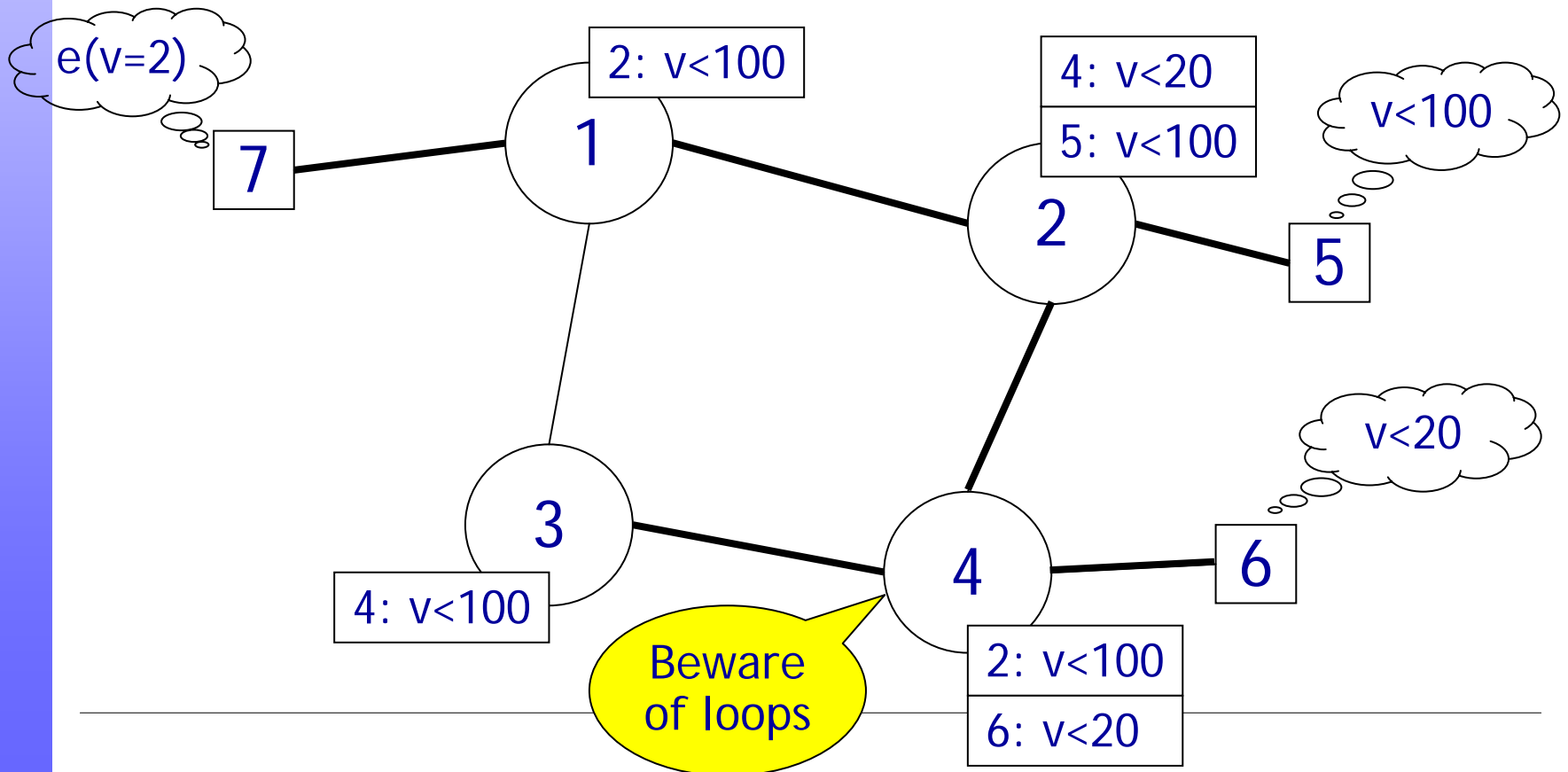
# Covered Subscriptions

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- Observations
  - Flooding the complete network for each subscription
  - Each broker has one routing table entry per subscription
- Idea: exploit redundancy of subscriptions
  - E.g., events matching „ $v < 20$ “ also match „ $v < 100$ “
  - We say: „ $v < 100$ “ covers „ $v < 20$ “
- Modified route setup: When a broker receives  $s$  from  $n$ 
  - Do nothing if  $s$  has been received before
  - Do nothing if there is an entry  $(n, s_i)$  and  $s$  is covered  $s_i$
  - Create routing table entry  $(n, s)$
  - Send  $s$  to all neighboring brokers

# Example

How to handle  
Unsubscriptions?



# Subscription Merging

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- Goal: reduce size of the routing table
- Approach: Merge two routing table entries with same destination into one
  - Perfect:  $(1, v < 10), (1, v > 10) \rightarrow (1, v \neq 10)$
  - Imperfect:  $(1, v < 10), (1, v > 11) \rightarrow (1, v \neq 10)$

# Advertisements

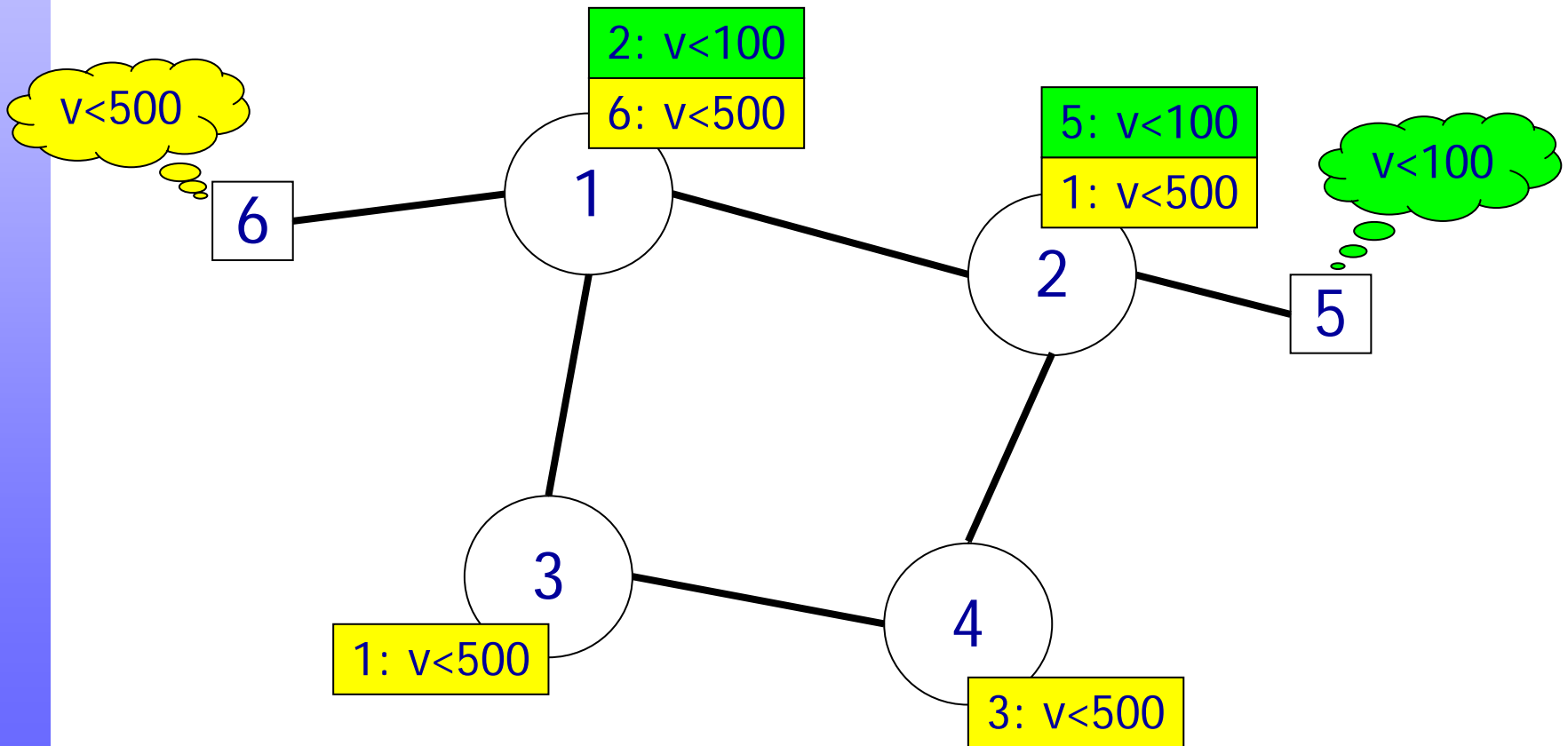
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- Idea: Publishers announce details on the events they will generate, e.g.  $v < 500$ 
  - Assuming this won't change over time
- Approach: Flood advertisements, building a second routing table
  - Subscriptions are no longer flooded, but routed using the advertisement routing table
  - Instead of checking if an event matches a subscription, need to check if a subscription overlaps an advertisement
    - E.g.,  $v < 500$  and  $v < 100$  overlap, but not  $v < 500$  and  $v > 600$

# Example

- Advertisement:  $v < 500$
- Subscription:  $v < 100$

What if advertisement after subscription?  
What if almost concurrently?





# Further Issues

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- Dynamic networks
  - Joining, leaving, crashing brokers
- Intermittent connectivity
  - Mobile networks with roaming publishers / subscribers
  - Event buffering
- Quality of Service
  - Event ordering, latency, guaranteed delivery, security, ...
- Composite events
  - A and B, A followed by B, ...